

1. Let  $C$  be a binary  $(n, k)$ -cyclic code with generator polynomial  $g(x) = (x + 1)p(x)$ , where  $p(x)$  does not divide  $x^t - 1$  for any  $t$ ,  $1 \leq t \leq n - 1$ . An error pattern of the form  $e(x) = x^i + x^{i+1}$  ( $0 \leq i \leq n - 1$ ) is called a *double-adjacent* error pattern.
  - (a) Prove that no two double-adjacent error patterns can be in the same coset of a standard array for  $C$ .
  - (b) Prove that  $C$  corrects all single errors and all double-adjacent errors.
  - (c) Construct a generator polynomial for a binary  $(15, 10)$ -cyclic code which corrects all single errors and all double-adjacent error patterns.
2. Consider  $GF(3^3)$  generated by  $f(x) = 1 - x + x^3$  over  $GF(3)$ .
  - (a) Prove that  $\beta = x^2$  has order 13 in  $GF(3^3)$ .
  - (b) Determine the cyclotomic cosets of 3 modulo 13.
  - (c) Use (a) and (b) to factor  $y^{13} - 1$  into irreducible monic polynomials over  $GF(3)$ .
3. Consider the finite field  $GF(2^5) = \mathbb{Z}_2[x]/(x^5 + x^2 + 1)$ . Then  $\alpha = x$  is a generator of  $GF(2^5)$ . We have the following minimal polynomials:

$$\begin{array}{ll}
 m_0(y) = y & m_{\alpha^5}(y) = 1 + y + y^2 + y^4 + y^5 \\
 m_1(y) = 1 + y & m_{\alpha^7}(y) = 1 + y + y^2 + y^3 + y^5 \\
 m_{\alpha}(y) = 1 + y^2 + y^5 & m_{\alpha^{11}}(y) = 1 + y + y^3 + y^4 + y^5 \\
 m_{\alpha^3}(y) = 1 + y^2 + y^3 + y^4 + y^5 & m_{\alpha^{15}}(y) = 1 + y^3 + y^5.
 \end{array}$$

Construct a generator polynomial for a binary  $(31, 11)$ -cyclic code which has designed distance 11.

4. Recall that  $C_{15}$  is a  $(15, 7)$ -BCH code over  $GF(2)$  with designed distance 5 generated by  $g(x) = m_{\beta}(x)m_{\beta^3}(x)$ , where  $\beta = y$  is an element of order 15 in  $GF(2^4) = \mathbb{Z}_2[y]/(y^4 + y + 1)$ . Decode the received word  $r = (00101\ 11000\ 11101)$  using  $C_{15}$ .

Solutions will be available on the course web site on April 2.